

Our Mission: Hacking Anything to Secure Everything

# Using Symbolic Execution for IoT Bug Hunting

Presenters: Grzegorz Wypych, X-Force Red

## Bio



Age - 36

Full name - Grzegorz Wypych (h0rac)

- Career path: Network Engineer => Network Architect => Software Developer => Security Researcher
- Languages: C, python, node.js, javascript, Java
- Papers: Yeah CCIE R&S will expire in 25/06/2019 together with other Cisco certs ;]
- Overall 15 years IT experience
- ARM/MIPS assembly enthusiast :)
- 0day CVEs on account related to TP-LINK devices
- When I do no research: I build fishing rods and fish out of the water :)
- I wish my day to have more than 24 hours :)
- Motto ? Before use.. disassemble :)



https://twitter.com/horac341



https://github.com/h0rac/



## 1) Problems with traditional dynamic vulnerability research

- You test where you are and not where you want to be
- Anti-debuggers applied
- Busybox without tftp, can't upload gdbserver, or core dump not available
- No ssh, no shell, no access (JTAG, UART)
- Qemu emulation nightmare

## You test where you are and not where you want to be

[ Legend: Modified register | Code | Heap | Stack | String | Szero: 0x0 \$at : 0x80353a0c \$v1 : 0x74 \$a0 : 0x7f81b7f6 → 0x00000a00 \$a1 : 0x00413a62 → 0x4b4f0000 \$a2 : 0x0 \$a3 : 0x0 \$t1 : 0x30 \$t3 : 0x7f81b439 → 0x00000031 ("1"?) \$t4 : 0x7f81b430 → "172.16.0.1" \$s0 : 0xffffffff \$s1 : 0x0042dd84 → 0x0042d198 → 0x0042dd84 → [loop detected] \$s3 : 0x00410000 → 0x00002021 ("!"?) \$55 : 0x00413a24 → \*JSESSIONID\* \$56 : 0x00413c10 → 0x00408330 → <http parser main+3568> move s0, zero \$57 : 0x0042d0a4 → 0x00413b98 → "Authorization" \$sp : 0x7f81b4f0 → 0x00000000 Sfir : 0x0 \$gp : 0x00435630 → 0x00000000 6x7181b4f0 +0x6008: 6x00600800 ← \$sp 0x7f81b4f4 +0x0004: 0x0000000 0x7f81b4f8 +0x0008: 0x0000000 8x7181b41c +0x600c: 0x00600000 0x7f81b500 +0x0010: 0x00435630 → 0x00000000 8x7f81b504 +0x0014: 0x0000000 8x7f81b508 +0x0018: 0x0000000 8x7f81b50c +0x001c: 0x00000000 → 0x408140 <http parser main+3072> lw gp, 16(sp) 0x408144 <http parser main+3076> bnez v0, 0x408178 <http parser main+3128> 0x408148 <http parser main+3080> move a0, s8 0x40814c <http parser main+3084> lw t9, -31724(gp) 0x408150 <http parser main+3088> nop 0x408154 <http parser main+3092> jalr [#0] Id 1, Name: "httpd", stopped, reason: SINGLE STEP [#0] 0x408140 → http parser main() [#1] 0x406040 → http inetd main() [#2] 0x404268 → http init main() [#3] 0x4034c0 → main() 0x00408140 in http parser main () get>

0x00408138 in http parser main ()

#### Traditional Debugging with GDB

- Hard to identify proper breakpoint place
- You follow single path
- Changing path selection with register modify could break execution
- Once debug fail, you need to start from the beginning.
- If this is remote debug session, you loose your breakpoints (Agh!!!)
- Time consuming !

But you can see step by step what is going on in process memory

#### Anti-debuggers applied

0042bcb8	int32_t (*	const	<pre>cmem_updateFirmwareBufFree@GOT)() = cmem_updateFirmwareBufFree</pre>
0042bcbc	int32_t (*	const	rdp_oidToOidStr@GOT)() = rdp_oidToOidStr
0042bcc0		const	signal@GOT)() = signal
0042bcc4	int32_t (*	const	dm_compareNumStack@GOT)() = dm_compareNumStack
0042bcc8	int32_t (*	const	dm_validateString@GOT)() = dm_validateString

0042bd7c	int32_t (* const sscanf@GOT)() = sscanf
0042bd80	<pre>int32_t (* const sigaction@GOT)() = sigaction</pre>
0042bd84	<pre>int32_t (* const setsid@GOT)() = setsid</pre>
0042bd88	int32_t (* const g_oidStringTable@GOT)() = g_oidStringTable
an internet	

When you try to hit breakpoint in debugger and step over or continue, instead going to expected destination you land in **SIGTRAP** ;/

To avoid, you can:

a) Try to patch binary, but no guarantee If this will work

- b) Try to use GDB for software debug bypass
- c) Modify registers to not execute during dynamic debugging

And then: We lost debug session... UPS :(

But we are missing our goal ! We don't want to spend time on avoiding anti-debuggers We want to utilize our time for vulnerability research and exploitation :)



Busy box without tftp, can't upload gdbserver or core dump not available

Standard binaries available under IoT OS usually have binding to busy box. First step After image retrieve is to check available commands under busybox. If we are lucky enough And tftp/ftp is available we have option to upload gdbserver binary for dynamic analysis, however Sometimes busybox is intentionally limited - What we can do then ?

We can try to reflash image with own busybox, but no guarantee it will work

If we want to have core dump for analysis, usually this commands enable it on device

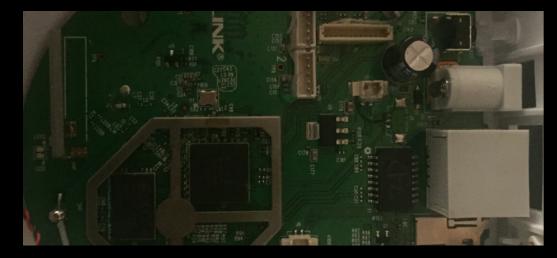
ulimit -c unlimited
echo /var/tmp/core > /proc/sys/kernel/core\_pattern

In some cases this do not work and you cannot grab core dumps for analysis



No ssh, no shell, no access

- Sometimes ssh is available but not for us :) What usually happens (example on TP-LINK devices), ssh is available only for certain application like Tether mobile app which is used for remote management.
- Telnet is usually limited to "cli" binary which is loaded in runtime. Every time user log in to device via telnet, binary is loaded to memory and provides limited config options.
- UART/JTAG not available, like on this IP Camera NC450



TP-LINK NC-450 board no JTAG/UART visible pins



Bonus - CLI binary TP-LINK devices Just as research bonus :) - I found something interesting in "cli" binaries available in all TP-LINK devices I've researched. It has hidden menu with shell access, but to enable it, it is required To have active debug session and manipulate "flags" in memory:

```
set {int}0x41e9d8 = 0x14 - g_cli_user_level
set {int}0x0041e9dc = 0x14 - g_cli_mode
```

#### This is normal command-line tool available on TP-LINK devices

elcome To Use TP-Li	nk COMMANE	)-LINE Interface Model.		
P-Link(conf)#help				
ormal mode commands				
clear		clear screen		
exit		leave to the privious mode		
help		help info		
history		show histroy commands		
logout		logout cli model		
onfig mode commands				
config		enter config mode		
igmp		igmp config		
wlctl		wireless config		
lan		lan config		
dev		device control		
usb		usb config		

## Bonus - CLI binary TP-LINK devices

#### This is how it looks like after modifying flags in memory

TP-Link(conf)#			
TP-Link?help			
normal mode command	s :		
clear		clear screen	
exit		leave to the privious mode	
help		help info	
history		show histroy commands	
logout		logout cli model	
privilege mode comm			
		enter privilege mode	
sh		force to cli	
config mode command	s :		
config		enter config mode	
igmp		igmp config	
wlctl		wireless config	
lan		lan config	
dev		device control	
usb		usb config	
TP-Link?sh			
[ doFshell ] cmd: s			
~ # ls -la			
drwxr-xr-x 10	138 web		
drwxr-xr-x 15	0 var		
drwxr-xr-x 4	38 usr		
dr-xr-xr-x 11	0 sys		
drwxr-xr-x 2	276 sbin		
dr-xr-xr-x 90	0 proc		
drwxr-xr-x 2	3 mnt		
lrwxrwxrwx 1	11 linux	xrc -> bin/busybox	
	1138 lib		
	502 etc		
drwxr-xr-x 7	1326 dev		
drwxr-xr-x 7 drwxr-xr-x 8	1326 dev 388 bin		
drwxr-xr-x 7 drwxr-xr-x 8 drwxr-xr-x 2			

I don't know if they left dev code for debug purposes or smth but why it is under production code in every device ? :)

#### Qemu emulation nightmare

Everyone is saying Qemu can emulate IoT binaries/firmware, let's verify that against real software :)

 TP-LINK devices usually store in flash memory "shared region" where they store configuration options. During Qemu emulation we do not have access to and strace immediately inform us about that and fail emulation.

sudo chroot . ./qemu-mipsel-static -strace usr/bin/httpd

```
40411 ipc(23,1234,0,950) = -i errno=22 (Invalid argument)
40411 write(1,0x7630f278,92)[ dm_shmInit ] 086: shmget to exitst shared memory failed. Could not create shared memory.
= 92
40411 ipc(1,-1,1,0) = -1 errno=22 (Invalid argument)
40411 write(1,0x7630f278,53)[ dm_acquireLock ] 252: lock failed, errno=22 rc=-1
= 53
qemu: uncaught target signal 11 (Segmentation fault) - core dumped
(angr) → rootfs
(angr) → rootfs
(angr) → rootfs sudo chroot . ./qemu-mipsel-static -strace usr/bin/httpd
```

00101021	11010		
00404328	lw	\$t9, -0x7d14(\$gp) {dm_shmInit@GOT}	
0040432c	nop		
00404330	jalr	\$t9	
00404334	move	\$a0, \$zero {0x0}	
00404338	jal	sub_403fb8	



## 2) Write plugin for Binary Ninja

Before we jump to symbolic execution, let's talk about Binary Ninja Disassembler

Pros:

- Nice python api
- Cheaper than IDA Pro
- Support multi-processors (ARM/MIPS/PowerPC etc)
- Modern UI :)
- Multi-disassembler options: Medium IL, ILL etc

Cons:

- Less features than IDA Pro
- Less processor types support
- No C decompiler

I think creators of Binary Ninja provides standard functionality to Disassembler, but leave a lot for users to add as plugins, and this is where power is unlimited



## Important Binary Ninja components

- When binary is loaded bv reference is available for us
- BinaryView and Architecture class allows to take basic information from analyzed binary (architecture, endianness, functions and their params etc)
- Most common utilized modules:
- A) plugin provides core for UI (PluginCommand, BackgroundTaskThread)
- B) interaction provides different UI components
- C) highlight colors for graph view

Python Console	*****
>>>	
>>>	
>>>	
>>> by	
<pre><binaryview: '="" 0x400000,="" bin="" firmware="" fmk="" home="" horac="" httpd',="" research="" rootfs="" start="" usr="" wr941nd="">&gt;&gt; bv.arch.name 'mips32' &gt;&gt;&gt; bv.get_function_at(0x4703f0) <func: mips32@0x4703f0=""> &gt;&gt;&gt; func = bv.get_function_at(0x4703f0) &gt;&gt;&gt; func.parameter_vars [<var arg1="" char*="">]</var></func:></binaryview:></pre>	len 0x1c7e00>
>>>	
Log Python Console	



## Plugin module (PluginCommand class)

There are two ways to use PluginCommand class from plugin module

### Use direct PluginCommand class in main python file

```
PluginCommand.register(
```

"Explorer\WR941ND\Explore", "Description", BackgroundTaskManager.vuln\_explore)

Encapsulate in separate class by inheritance

class UIPlugin(PluginCommand):

```
def __init__(self):
    super(UIPlugin, self).register_for_address("Explorer\WR941ND\Start Address\Set",
    "Set execution starting point address", self.set_start_address)
    super(UIPlugin, self).register("Explorer\WR941ND\Start Address\Clear",
```

Explanation on function and parameters:

register - expect handler with one param, bv instance
register\_for\_address - expect handler with two params, bv instance and address

"\" is important it allows to create sub-menus



## Plugin module (BackgroundThread class)

To execute actions in BinaryNinja, we need to inherit from **BackgroundTaskThread** And override **run** method by our implementation

self.explorer.run()

```
class AngrRunner(BackgroundTaskThread):
    def __init__(self, bv, explorer):
        BackgroundTaskThread.__init__(
            self, "Vulnerability research with angr started...", can_cancel=True)
        self.bv = bv
        self.explorer = explorer
    def run(self):
```

We can define own parameters for \_\_init\_\_ constructor. Here we provide own explorer instance which in this example could be VulnerabilityExplorer, ROPExplorer, JSONExploitCreator, FileExploitCreator



# Interaction module

UI components are provided by interaction module. They are very easy to use

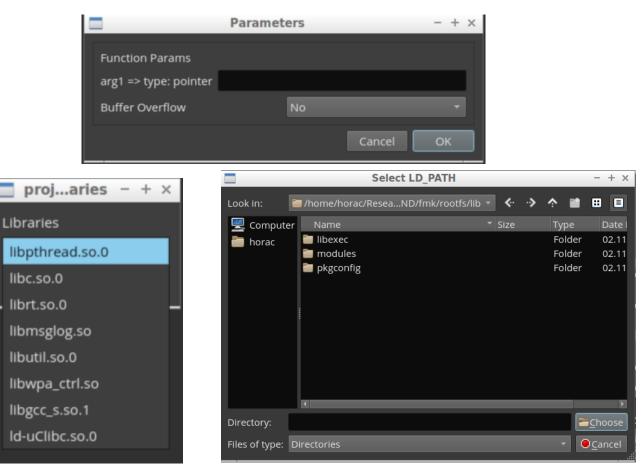
```
menu_items = self.generate_menu_text_fields(mapped_types)
menu = interaction.get_form_input(menu_items, "Parameters")
```

We can create UI components separately or use function **get\_form\_input** to create our custom menu. Function expect list of fields we want to include (TextLineField, ChoiceField etc). It returns also list of results



## Interaction module (sample UI)

Sample UI look of components.





# Highlight module

#### @classmethod

```
def color_path(self, bv, addr):
    # Highlight the instruction in green
    blocks = bv.get_basic_blocks_at(addr)
    UIPlugin.path.append(addr)
    for block in blocks:
        block.set_auto_highlight(HighlightColor(
            HighlightStandardColor.GreenHighlightColor, alpha=128))
        block.function.set_auto_instr_highlight(
            addr, HighlightStandardColor.GreenHighlightColor)
```

This is example function used for path coloring during symbolic execution. We first get basic blocks Of assembly by address and highlight them to whatever color we want. Later also single addresses are colored. Results are store in class variable path.

We can call this function from any place, but in plugin I will present I use it during symbolic execution

# 3) How you can search for vulnerabilities without hacking physical device access or without Qemu emulation

angr features we will use

- We will look on CVE -2019-6989 Buffer Overflow WR941ND (MIPS)
- We will identify vulnerable code with basic static analysis
- We will confirm vulnerability with symbolic execution (well.. tuned a little :)) using created

plugin

And guess how ? We will not even try to run firmware, we will emulate it with angr



## angr features

#### Load/Save to emulated memory

```
state.memory.store(sp+0x2c, state.solver.BVV(self.gadget3, 32))
state.memory.load(0x100, size)
```

#### Load/Save to register

```
state.regs.s0 = 0x100
state.regs.s1 = "AAAA"
pc = state.solver.eval(state.regs.pc, cast_to=int)
s1 = state.solver.eval(state.regs.s1, cast to=bytes)
```

#### **CFG Analysis**

self.proj.analyses.CFGFast(regions=[(self.func\_start\_addr, self.func\_end\_addr)])

#### Hooking

```
self.proj.hook(self.func_end_addr, self.overwrite_ra)
```

#### **PointerWrapper**

```
angr.PointerWrapper(item.get('value'))
```

#### Call state

self.proj.factory.call\_state(self.func\_start\_addr, args['arg0'])



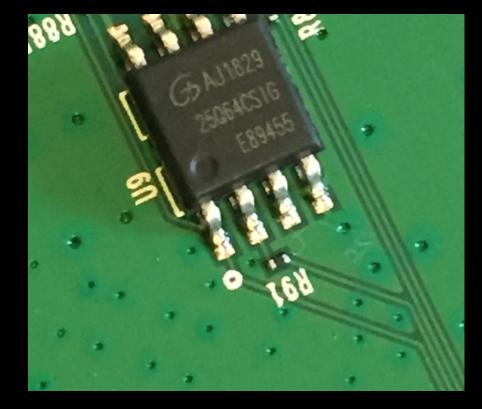
## Find vulnerable Endpoint-W941ND

In Management panel, we have option to send health pings and check availability. However **m**odyfing **ping\_addr** with custom string crash httpd service in router. Now we know something is wrong but what exactly ???

Go Cantel < Y > Y	Target: http://192.168.0.1
Request	Response
Raw         Params         Headers         Hex           GET         /ml_ZTGSWBORNIIIMA/userRpm/PinglframeRpm.htm?ping_addr=AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA	Raw Headers Hex HTTP/1.1 200 CK Server: Router Webserver Connection: close Content Type: text/html WWW-Authenticate: Basic realm="TP-LINK Wireless N Router WR940N"

Let's dump firmware and start some basic static analysis and search for strings like URL endpoint or parameters names

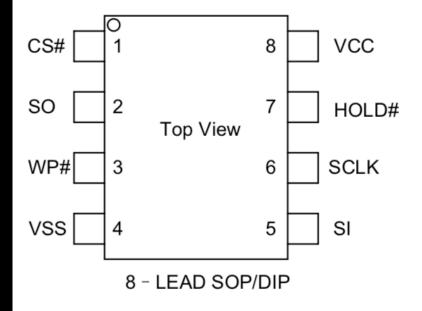




Flash chip (GD25Q64C) on most TP-LINK devices - Archer C5 v4 - another device but same process

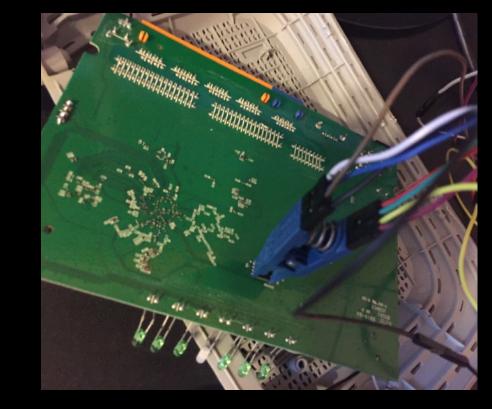


## **Connection Diagram**



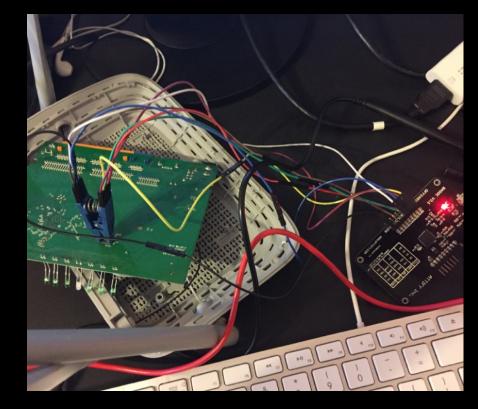
Chip info how to connect PINS for SPI



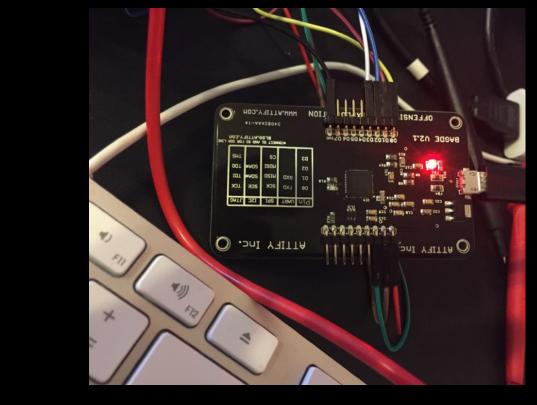


SOC8 clips connected to flash chip





Clips connected to flash chip and Attify Badge, bus pirate and any other SPI supported device will also work



Connected Attify badge over SPI sudo ./flashrom -p ft2232\_spi:type=232H -r firmware.bin



## No device, but firmware available on vendor site

#### https://www.tp-link.com/us/support/download/tl-wr941nd/#Firmware

TL-WR941ND(US)_V6_151203 ≚				
Published Date: 2016-12-03	Language: English	File Size: 3.21 MB		
Notes: TL-WR940N(US)3.0 / TL-WR941ND(US)6.0				

Next step is to extract firmware using firmware-mod-kit (easiest way) and find binaries in rootfs we want to analyse.

~/Research/firmware-mod-kit/extract-firmware.sh wr941nd.bin

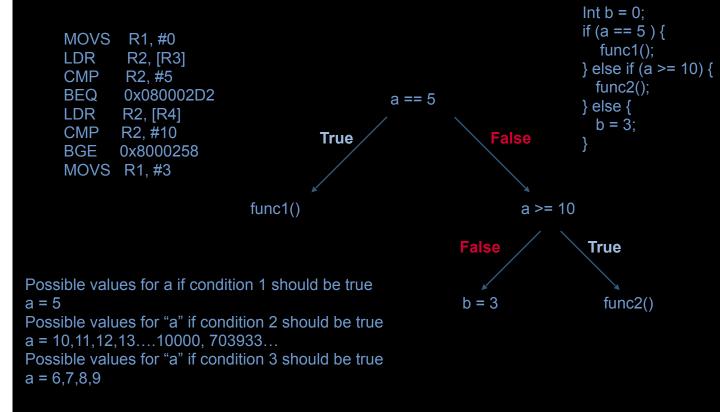


## Firmware - first look

plugins cd ~/Research/firmware/WR941ND WR941ND ls WR941ND cd fmk/rootfs rootfs ls bin dev etc lib linuxrc mnt proc gemu-mips-static root sbin sys 🎹 usr var web rootfs cd usr/bin bin ls arping dbclient dropbear dropbearconvert dropbearkey httpd lld2d logger scp test tftp bin bin bin pwd /home/horac/Research/firmware/WR941ND/fmk/rootfs/usr/bin bin readelf -h httpd ELF Header: Magic: 7f 45 4c 46 01 02 01 00 00 00 00 00 00 00 00 00 Class: ELF32 2's complement, big endian Data: 1 (current) Version: OS/ABI: UNIX - System V ABI Version: 0 Type: EXEC (Executable file) Machine: MIPS R3000 Version:  $0 \times 1$ Entry point address: 0x41c5b0 Start of program headers: 52 (bytes into file) Start of section headers: 0 (bytes into file) 0x70001007, noreorder, pic, cpic, o32, mips32r2 Flags: Size of this header: 52 (bytes) Size of program headers: 32 (bytes) Number of program headers: 9 Size of section headers: 0 (bytes) Number of section headers: Section header string table index: 0 bin readelf -d httpd Dvnamic section at offset 0x180 contains 28 entries: Type Name/Value Taq 0x00000001 (NEEDED) Shared library: [libpthread.so.0] Shared library: [libc.so.0] Shared library: [librt.so.0] Shared library: [libmsglog.so] Shared library: [libutil.so.0] Shared library: [libwpa\_ctrl.so] 0x00000001 (NEEDED) 0x00000001 (NEEDED) 0x00000001 (NEEDED) 0x00000001 (NEEDED) 0x00000001 (NEEDED) Shared library: [libgcc\_s.so.1] 0x00000001 (NEEDED) 0x0000000c (INIT) 0x41c524 0x0000000d (FINI) 0x543f30

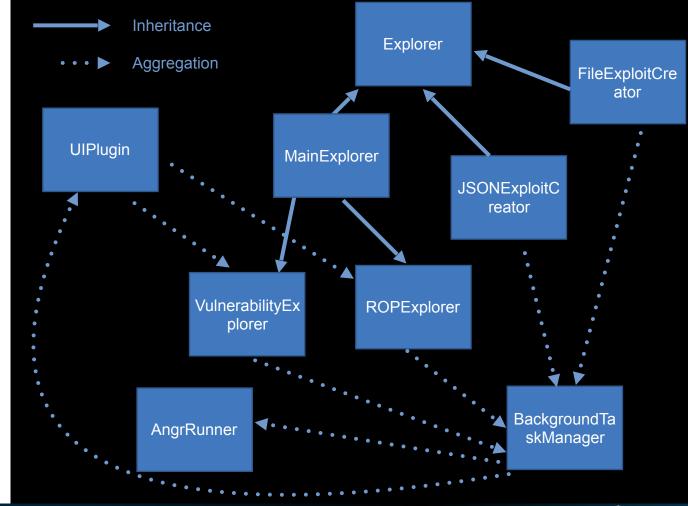


## Symbolic execution in nutshell



#### Example of symbolic execution tree base on ARM CPU instructions

## Explorer view





## Vulnerability exploration DEMO

## Vulnerability exploration DEMO TIME :)



4) I have RA in control, let's build PoC exploit without debugger /crash dump/memory snapshot ?

- MIPS Assembly in nutshell
- We will present features of angr we will use for ROP exploitation
- We will present gadgets for ROP
- We will create ROP chain
- We will execute ROP chain
- We will provide report for CPU registers and stack during ROP execution

And guess how ? We will not even try to run firmware, we will emulate it with angr





# MIPS Assembly in nutshell

• Endianness: Little Endian(MIPSEL) and Big Endian(MIPS)

- First four arguments to function passed in registers (\$a0, \$a1, \$a2, \$a3)
- Function need more arguments ? They are pushed on stack
- Instruction Pointer aka intel EIP => \$pc (Program counter)
- Stack pointer: \$sp
- Calling function executed by loading register to \$t9 and jalr \$t9
- Return address stored in \$ra
- Return value \$v0
- Callee responsible to store value of registers before executing
- Space for local variables in stack frame: \$sp, sp,- 0x3c in prologue



## **ROP** gadgets

### Gadget 1

### 0x00055c60:

addiu \$a0, \$zero, 1; # prepare param for sleep func move \$t9, \$s1; # copy gadget2 address jalr \$t9;

### Gadget 2

### 0x00024ecc:

lw \$ra, 0x2c(\$sp); # load gadget 3 address lw \$s1, 0x28(\$sp); # load sleep func addr lw \$s0, 0x24(\$sp); # load junk jr \$ra;

### Gadget 3

#### 0x0001e20c:

move \$t9, \$s1; lw \$ra, 0x24(\$sp); # load gadget 4 address lw \$s2, 0x20(\$sp); # load junk lw \$s1, 0x1c(\$sp); # load gadget 5 address lw \$s0, 0x18(\$sp); # load junk jr \$t9;



## **ROP** gadgets

### Gadget 4

#### 0x000195f4:

addiu \$s0, \$sp, 0x24; # store in \$s0 address of shell code move \$a0, \$s0; # copy shell code address to \$a0 move \$t9, \$s1; # copy address of gadget5 to \$t9 jalr \$t9; # jump

### Gadget 5

#### 0x000154d8: move \$t9, \$s0; # copy address of \$s0 to \$t9 jalr \$t9; # execute shell code

Sleep function

0x00053ca0

## ROP exploitation DEMO

## ROP exploitation DEMO TIME :)





# THANK YOU

#### FOLLOW US ON:



securityintelligence.com

xforce.ibmcloud.com



youtube/user/ibmsecuritysolutions

© Copyright IBM Corporation 2018. All rights reserved. The information contained in these materials is provided for informational purposes only, and is provided AS IS without warranty of any kind, express or implied. Any statement of direction represents IBM's current intent, is subject to change or withdrawal, and represent only goals and objectives. IBM, the IBM logo, and other IBM products and services are trademarks of the International Business Machines Corporation, in the United States, other countries or both. Other company, product, or service names may be trademarks or service marks of others.

Statement of Good Security Practices: IT system security involves protecting systems and information through prevention, detection and response to improper access from within and outside your enterprise. Improper access can result in information being altered, destroyed, misappropriated or misused or can result in damage to or misuse of your systems, including for use in attacks on others. No IT system or product should be considered completely secure and no single product, service or security measure can be completely effective in preventing improper use or access. IBM systems, products and services are designed to be part of a lawful, comprehensive security approach, which will necessarily involve additional operational procedures, and may require other systems, products or services to be most effective. IBM does not warrant that any systems, products or services are immune from, or will make your enterprise immune from, the malicous or illegal conduct of any party.

